### Introduction to Objective Caml

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October 30, 2006

- Motivation
- Simple Expressions and Types
- Functions
- Simple Pattern Matching
- Compound Datatypes
- I/O and Compilation
- Resources

Why OCaml?

- Convenient encoding of tree structures
- Powerful pattern matching facilities
- Close correspondence to mathematical notation

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#### Features of OCaml

- Functional language (functions are first-class values)
- Strong and statically typed
- Parametric polymorphism
- Type inference
- Recursive, algebraic datatypes (trees, lists, ...)
- Garbage collection
- Modul-system
- (Object-system)

Motivati

# Applications written in OCaml

- File sharing: MLdonkey (http://mldonkey.org/)
- File synchronizer: unison
   (http://www.cis.upenn.edu/~bcpierce/unison/)
- Compilers and interpreters: OCaml, XQuery, XDuce, CDuce
- Proof assistant: Coq (http://coq.inria.fr/)

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Simple Expressions and Types

#### The Toplevel Loop

- Interactive development
- Evaluation of expressions (calculator)
- Definitions

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# Basic Types (1)

```
# ();;
- : unit = ()
```

- Singleton type: () is the only element of unit
- Similar to void in C or Java
- Result type of functions with side effects

Simple Expressions and Types

```
# 2 + 5 * 8;;
- : int = 42
```

- Signed integers, represented by a machine word minus one bit
- Common Operators: +, -, \*, /, mod

Simple Expressions and Types

• Conversions: string\_of\_int, int\_of\_string, float\_of\_int

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Simple Expressions and Types

#### Basic Types (2)

```
# 3.1415926536 *. 2.0;;
- : float = 6.2831853072
```

- IEEE double-precision floating point, equivalent to C's double
- Arithmetic operators end with a dot: +., -., \*., /.
- Conversions: string\_of\_float, int\_of\_float

```
# Char.uppercase 'x';;
- : char = 'X'
```

- Latin-1 characters (unicode library: http://camomile.sf.net/)
- Functions: Char.lowercase, Char.uppercase
- Conversions: Char.code (character → integer), Char.chr (integer → character)

Basic Types (3)

```
# "Hello " ^ "World\n";;
-: string = "Hello World\n"
```

- Strings with Latin-1 encoding
- Operators: ^ (concatenation), "Hello".[1] (index access)
- Functions: String.length, String.sub

```
# 1 = 2 || false;;
- : bool = false
```

- Operators: & &, | |, not
- Comparisons: = (equality), <> (inequality), <, <=, >, >= These operators work on arbitrary but equal types; for some types, a runtime exception is raised.

Simple Expressions and Types

#### Conditionals and Variables

#### Conditionals

```
# if 1 < 2 then 3 + 7
   else (if "Hello" = "stefan" then 0 else 42);;
- : int = 10</pre>
```

#### **Variables**

- Variables are names for values
- No assignment!

```
# let x = 4;;
val x : int = 4
# 38 + x;;
- : int = 42
# let y = 3 in 39 + y;;
- : int = 42
# y;;
Unbound value y
```

#### **Functions**

```
# let square x = x * x ;;
val square : int -> int = <fun>
# square 42;;
- : int = 1764
```

- Function type: t1 -> t2
- Function call without parenthesis around argument

```
# let average x y = (x + y) / 2;;
val average : int -> int -> int = <fun>
# average 21 63;;
- : int = 42
```

- Type of multi-argument functions: t1 -> t2 -> ... -> tn
- Function call: concatenate all arguments to the function

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**Functions** 

#### **Nested Functions**

• Functions may be arbitrarily nested.

```
# let sum_of_3 x y z =
    let sum a b = a + b
        in sum x (sum y z);;
val sum_of_3 : int -> int -> int -> int = <fun>
# sum_of_3 1 2 3;;
- : int = 6
```

Function

#### **Recursive Functions**

- A recursive function calls itself inside its own body.
- Defined as ordinary functions, but uses let rec instead of let.
- Example: function that computes  $x^i$

```
# let rec power i x =
    if i = 0 then
        1.0
    else
        x *. (power (i - 1) x);;
val power : int -> float -> float = <fun>
# power 5 2.0;;
- : float = 32.
```

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**Functions** 

#### Mutually Recursive Functions

• Connect several let rec definitions with the keyword and.

```
# let rec f i j =
    if i = 0 then
    j
    else g (j - 1)
and g j =
    if j mod 3 = 0 then
     j
    else f (j - 1) j;;
val f : int -> int -> int = <fun>
val g : int -> int = <fun>
# g 5;;
- : int = 3
```

Polymorphic Functions

- Work on values of arbitrary type
- Arbitrary types represented as type variables 'a, 'b, ...

```
# let id x = x;;
val id : 'a -> 'a = <fun>
```

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unctions

#### The Value Restriction

- Only values can be polymorphic.
- Function applications are not values.
- The value restriction is needed to ensure soundness in the presence of side-effects.

```
# let id' = id id;;
val id' : '_a -> '_a = <fun>
# id' 5;;
- : int = 5
# id';;
- : int -> int = <fun>
```

Functio

# Higher-order functions

- Functions are ordinary values.
- A higher-order function takes another function as an argument or returns it as the result.
- Partial application of a function (with less arguments than expected) returns another function

```
# let add x y = x + y;;
val add : int -> int -> int = <fun>
# let inc = add 1;;
val inc : int -> int = <fun>
# let compose f g x = f (g x);;
val compose : ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b = <fun>
# compose inc inc 0;;
- : int = 2
```

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**Functions** 

### Function Types in Detail

• The arrow associates to the right: The type

```
int -> int -> int is the same as
int -> (int -> int)
```

- add takes an int and returns a function of type int -> int
- Function application associates to the left: The expression add 1 2 is the same as (add 1) 2
- The sub-expression (add 1) has type int -> int so we can apply it to the integer 2

```
# let inc = add 1;;
val inc : int -> int = <fun>
# inc 2;;
- : int = 3
```

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**Anonymous Functions** 

• The keyword fun constructs an anonymous function.

```
# fun x -> x + 1;;
- : int -> int = <fun>
# compose inc (fun x -> x + 1) 0;;
- : int = 2
```

• Definitions such as let add x y = x + y are just syntactic sugar. Here is the expanded definition:

```
# let add = fun x y -> x + y;;
val add : int -> int -> int = <fun>
```

Simple Pattern Matching

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Simple Pattern Matching

#### Simple Pattern Matching

- Powerful feature
- Defines expressions by case analysis
- Simple pattern: constant or variable

Constant matches only the constant value given
Variable matches all values and binds the value to the
variable

```
# let rec fib i =
    match i with
        0 -> 0
        | 1 -> 1
        | j -> fib (j - 2) + fib (j - 1);;
val fib : int -> int = <fun>
# fib 1;;
- : int = 1
# fib 6;;
- : int = 8
```

## Matching Order

- Cases of a match expression are tried in sequence, from top to bottom.
- The body of the first matching case is evaluated.
- The following definition of fib is wrong (fib loops forever when called).

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Simple Pattern Matching

### **Incomplete Matches**

OCaml issues a warning if the cases of a match do not cover all possible values:

```
# let rec fib i =
    match i with
      0 -> 0
      | 1 -> 1;;
Warning P: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
2
val fib : int -> int = <fun>
# fib 2;;
Exception: Match_failure ("", 55, 2).
```

Simple Pattern Matching

# Functions with Matching

- Common situation: pattern matching on the last argument of a function
- Use the function keyword instead of an explicit match expression

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Simple Pattern Matching

## **Matching Characters**

#### With pattern ranges and wildcard pattern:

Simple Pattern Matching

## Matching Strings

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Simple Pattern Matching

## Patterns Everywhere

• Patterns are used in all binding mechanisms:

```
let pattern = expression
let name pattern ... pattern = expression
fun pattern -> expression
```

Very useful with tuples and records (introduced next)

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# Tuples

Fixed-length sequences of values with arbitrary types

Compound Datatypes

• Construction:

• Elimination by pattern matching:

```
# let (a, b, c) = p;;
val a : string = "2-times"
val b : int -> int = <fun>
val c : bool = true
# match p with (a, _, _) -> a;;
- : string = "2-times"
```

• Pairs can be eliminated with fst and snd:

```
# fst (1,2);;
- : int = 1
# snd (1,2);;
- : int = 2
```

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Compound Datatypes

#### Lists

- Variable-length sequences of values with the same type
- Two constructors:

```
Nil [], the empty list Cons e_1::e_2, creates a new list with first element e_1 and rest of the list e_2
```

Shorthand notation:

```
[e_1; \ldots; e_n] is identical to e_1:: (e_2:: \ldots:: (e_n:: []) \ldots)
```

 $\bullet$  t  $\,$  list is the type of lists with elements of type t

```
# let l = "Hello" :: "World" :: [];;
val l : string list = ["Hello"; "World"]
# let l' = [1;2;3];;
val l' : int list = [1; 2; 3]
```

Compound Datatype

## Lists and Pattern Matching

Lists are eliminated using pattern matching:

Compound Datatypes

## The Map Function

- The function List.map applies a function to every element in a list
- List.map : ('a -> 'b) -> 'a list -> 'b list
- We can define the function inc\_list in terms of map because

```
inc_list [i1; ...; in] = [i1+1; ...; in+1].
```

```
# let inc_list = List.map (fun i -> i+1);;
val inc_list : int list -> int list = <fun>
# inc_list [1; 2; 3; 4];;
- : int list = [2; 3; 4; 5]
```

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Compound Datatypes

#### The Fold Function

- The function List.fold\_right "folds" a function over a list
- List.fold\_right :
   ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b
- We can define the function sum\_list in terms of map because sum list [i1; ...; in] = i1 + ... + in.

```
# let sum_list 1 = List.fold_right (+) 1 0;;
val sum_list : int list -> int = <fun>
# sum_list [1; 2; 3; 4];;
- : int = 10
```

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Compound Datatypes

#### Algebraic datatypes

- They represent the union of several different types.
- Every alternative has an unique, explicit name.
- General syntax:

```
type typename =
   Name<sub>1</sub> of type<sub>1</sub>
| Name<sub>2</sub> of type<sub>2</sub>
| Name<sub>n</sub> of type<sub>n</sub>
```

- The names *Name<sub>i</sub>* are called constructors; they must start with a capital letter.
- The part of *type*; is optional.

Compound Datatypes

#### Example

Compound Datatypes

## Pattern Matching with Algebraic Datatypes

```
# let float_of_number = function
    Zero
    -> 0.0
    | Integer i
     -> float_of_int i
    | Fraction (i,j)
     -> float_of_int i /. float_of_int j;;
val float_of_number : number -> float = <fun>
# float_of_number semi;;
- : float = 0.5
```

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# Binary Trees

```
# type 'a tree = Node of ('a * 'a tree * 'a tree) | Leaf;;
type 'a tree = Node of ('a * 'a tree * 'a tree) | Leaf
\# let rec insert x = function
      Leaf -> Node (x, Leaf, Leaf)
    | Node (y, l, r) ->
       if x < y
           then Node (y, insert x l, r)
           else if x > y then Node (y, l, insert x r)
           else Node (y, l, r);;
val insert : 'a -> 'a tree -> 'a tree = <fun>
# let tree = Node (5, Node (1, Leaf, Leaf),
                     Node (7, Leaf, Leaf));;
val tree : int tree = Node (5, Node (1, Leaf, Leaf),
                              Node (7, Leaf, Leaf))
# let tree' = insert 6 tree;;
val tree' : int tree =
 Node (5, Node (1, Leaf, Leaf),
           Node (7, Node (6, Leaf, Leaf), Leaf))
```

Compound Datatypes

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#### Compound Datatypes

### The Option Type

- Important builtin type
- Used to write partial functions

```
# type 'a option = Some of 'a | None;;
type 'a option = Some of 'a | None
```

#### Records

• Labeled collections of values with arbitrary types

Compound Datatypes

Record types must be declared

```
# type point = { point_x : int; point_y : int};;
type point = { point_x : int; point_y : int; }
```

- Label names must be globally unique
- Record construction:

```
# let p = { point_x = 5; point_y = 3 };;
val p : point = {point_x = 5; point_y = 3}
```

Field selection:

```
# let move p1 p2 =
    { point_x = p1.point_x + p2.point_x;
        point_y = p1.point_y + p2.point_y };;
val move : point -> point -> point = <fun>
# move p p;;
- : point = {point_x = 10; point_y = 6}
```

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I/O and Compilation

IO

#### Some functions for doing I/O:

```
val print_string : string -> unit
val print_endline : string -> unit
val prerr_string : string -> unit
val prerr_endline : string -> unit
val read_line : unit -> string

val open_out : string -> out_channel
val output_string : out_channel -> string -> unit

val open_in : string -> in_channel
val input_line : in_channel -> string
```

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I/O and Compilation

## Compilation

- Files with OCaml source code have the extension .ml
- Compiler ocamle: produces portable bytecode
- Compiler ocamlopt: produces fast native code
- Compiled program executes definitions in order of their appearance in the source file(s)
- Functions in some other source file foo.ml must be qualified with the prefix Foo.
- If file bar.ml uses functions from foo.ml, then bar.ml must come after foo.ml on the commandline. No cycles are allowed!

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I/O and Compilation

#### Compilation Example

• File fib.ml:

```
let fib = \dots
```

• File main.ml:

```
let _ =
  let _ = print_string "Input some number: " in
  let line = read_line () in
  let i = int_of_string line in
  let j = Fib.fib i in
    print_endline ("Result: " ^ string_of_int j)
```

- Compilation: ocamlc -o fib fib.ml main.ml
- Produces file fib:

```
$ ./fib 6
Input some number: 6
Result: 8
```

Resources

- OCaml Homepage: http://caml.inria.fr/
- Language Manual: http:

//caml.inria.fr/pub/docs/manual-ocaml/index.html

The standard library:

http://caml.inria.fr/pub/docs/manual-ocaml/libref/

- Jason Hickey: Introduction to the Objective Caml Programming Language. (The slides are based on this script)
   http://files.metaprl.org/doc/ocaml-book.pdf
- Emmanuel Chailloux, Pascal Manoury and Bruno Pagano:
   Developing Applications with Objective Caml

http://caml.inria.fr/pub/docs/oreilly-book/html /index.html

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