

# Softwaretechnik

## Lecture 08: Testing and Debugging — Overview

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## Essential Reading

- ▶ *Why Programs Fail: A Guide to Systematic Debugging*, A Zeller
- ▶ *The Art of Software Testing*, 2nd Edition, G J Myers

## Further Reading

- ▶ *Code Complete*, 2nd Edition, S McConnell

**\$ 60 billion**

yearly cost of software errors for US economy [NIST 2002]

**\$ 180 billion**

total sales of software in 2000

697,000 software engineers & 585,000 computer programmers

# Cost of Software Errors

estimated

**50%**

of each software project spent on testing  
(spans from 30% to 80%)

# Cost of Software Errors

very rough approximation

money  
spent on  
testing  $\approx$  cost of  
remaining  
errors

# Cost of Software Errors

very rough approximation

$$\begin{array}{rcc} \text{money} & & \text{cost of} \\ \text{spent on} & + & \text{remaining} \\ \text{testing} & & \text{errors} \\ & = & \end{array}$$

66% of size of software  
industry

# A Quiz About Testing

A simple program

## Input

Read three integer values from the command line.

The three values represent the lengths of the sides of a triangle.

## Output

Tells whether the triangle is

**Scalene:** no two sides are equal

**Isosceles:** exactly two sides are equal

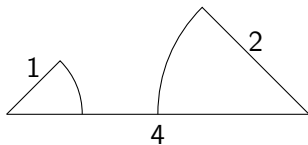
**Equilateral:** all sides are equal

Task: Create a Set of **Test Cases** for this Program



## Solution — 1 Point for each Correct Answer

Q 1: (4,1,2) a **invalid** triangle



Why not a valid triangle?  $(a,b,c)$  with  $a > b + c$

Define valid triangles:  $a \leq b + c$

## Solution — 1 Point for each Correct Answer

Q 2: some permutations of previous (1,2,4), (2,1,4)

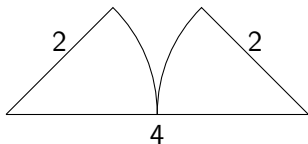
Fulfill above definition, but are still invalid.

Patch definition of valid triangles:

$$a \leq b + c \text{ and } b \leq a + c \text{ and } c \leq a + b$$

## Solution — 1 Point for each Correct Answer

Q 3: (4,2,2) a **invalid** triangle with **equal** sum



Fulfills above definition, but is invalid (depending on what we want!).

Patch definition of valid triangles:

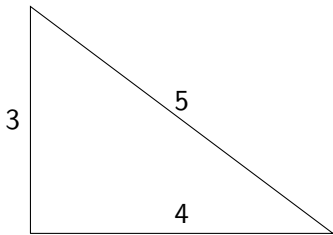
$$a < b + c \text{ and } b < a + c \text{ and } c < a + b$$

## Solution — 1 Point for each Correct Answer

Q 4: some permutations of previous  $(2,2,4)$ ,  $(2,4,2)$

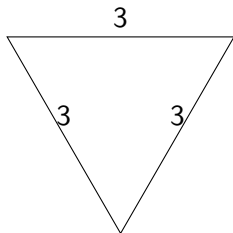
## Solution — 1 Point for each Correct Answer

Q 5: (3,4,5) a **valid scalene** triangle



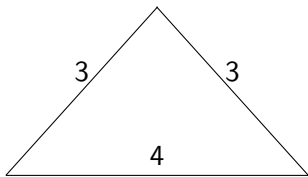
## Solution — 1 Point for each Correct Answer

Q 6: (3,3,3) an **equilateral** triangle



## Solution — 1 Point for each Correct Answer

Q 7: (3,4,3) **valid isosceles** t.



## Solution — 1 Point for each Correct Answer

Q 8: all permutations of valid isosceles triangle:

$(3,4,3)$ ,  $(3,3,4)$ ,  $(4,3,3)$



## Solution — 1 Point for each Correct Answer

Q 9: one side with **zero** value (0,4,3)

## Solution — 1 Point for each Correct Answer

Q 10: one side with **negative** value  $(-1,4,3)$

## Solution — 1 Point for each Correct Answer

Q 11: all sides zero (0,0,0)

## Solution — 1 Point for each Correct Answer

Q 12: at least one value is non-integer (1,3,2.5)

## Solution — 1 Point for each Correct Answer

Q 13: wrong number of arguments (2,4) or (1,2,3,3)

## Solution — 1 Point for each Correct Answer

Q 14 (the most important one):

Did you specify the expected output in each case?

## About the Quiz

- ▶ Q 1–13 correspond to failures that have actually occurred in implementations of the program
- ▶ How many questions did you answer?  
< 5? 5 – 7? 8 – 10? > 10? All?
- ▶ Highly qualified, experienced programmers score **7.8** on average

# First Conclusions

- ▶ Finding good and sufficiently many test cases is difficult
- ▶ Even a good set of test cases cannot exclude more failures
- ▶ A **specification** is required to identify failures

The discipline of Testing is all about Test Cases

Remark: At Ericsson: 35% of code is test cases!





# Failure and Specification

Some failures are obvious

- ▶ obviously wrong output/behaviour
- ▶ non-termination
- ▶ crash
- ▶ freeze

... but most are not!

In general, what constitutes a failure, is defined by a **specification!**

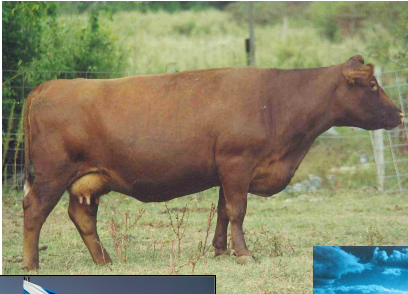
**Correctness is a relative notion**

— B. Meyer, 1997

**Every program is correct with respect to SOME specification**

— myself, today

# Specification: Intro



## Economist:

The cows in Scotland are brown

## Logician:

No, there are cows in Scotland of which one at least is brown!

## Computer Scientist:

No, there is at least one cow in Scotland, which is brown on one side!!

# Specification: Putting it into Practice

## Example

A Sorting Program:

```
1 public static Integer[] sort(Integer[] a) { ...  
    }
```

Testing `sort()`:

- ▶ `sort({3, 2, 5}) == {2, 3, 5}` ✓
- ▶ `sort({}) == {}` ✓
- ▶ `sort({17}) == {17}` ✓

Specification

*Requires:* *a* is an array of integers

*Ensures:* returns the sorted argument array *a*

## Example Cont'd

### Example

```
1 public static Integer[] sort(Integer[] a) { ...  
    }
```

### Specification

*Requires:* a is a *non-null* array of integers

*Ensures:* returns the *unchanged* reference a containing a permutation of the *old* contents of a that is sorted

sort(null) throws NullPointerException ❌

# The Contract Metaphor

**Contract** is preferred specification metaphor for procedural and OO PLs

first propagated by B. Meyer, *Computer* 25(10)40–51, 1992

Same Principles as Legal Contract between a Client and Supplier

**Supplier** aka **Implementer**, in JAVA, a class or method

**Client** Mostly a caller object, or human user for `main()`

**Contract** One or more pairs of **ensures/requires** clauses defining mutual benefits and obligations of client and implementer

# The Meaning of a Contract

Specification (of method  $C::m()$ )

*Requires:* Precondition

*Ensures:* Postcondition

*"If a caller of  $C::m()$  fulfills the **required Precondition**, then the class  $C$  **ensures** that the **Postcondition** holds after  $m()$  finishes."*

Often the following **wrong** interpretations of contracts are seen:

**Wrong!**

*"Any caller of  $C::m()$  must fulfill the **required Precondition**."*

**Wrong!**

*"Whenever the **required Precondition** holds, then  $C::m()$  is executed."*

# Failure

Definition: **failure**

A method **fails** if it is called in a state fulfilling the required precondition of its contract and does not terminate in a state fulfilling the postcondition.

Non-termination, abnormal termination considered as failures here



# Notions of Correctness

Definition: **partial correctness**

A method is **partially correct** if whenever it is started in a state fulfilling the required precondition and it terminates, then its final state fulfills the postcondition.

This amounts to proving **Absence of Failures!**

Definition: **total correctness**

A method is **totally correct** if whenever it is started in a state fulfilling the required precondition, then it terminates and its final state fulfills the postcondition.

Total correctness implies termination!

# Invariant

Objects with non-trivial state  
often maintain a **class invariant**.

Example: a class for dates

```
public class Date {  
    public int day;  
    public int month;  
    public int year;  
}
```

Invariant:

```
1 <= day <= 31 /\ 1 <= month <= 12 /\  
(month in {4, 6, 9, 11} => day <= 30) /\  
(month == 2 => day <= 29) /\  
(month == 2 /\ (year % 4 != 0 \/ (year % 100 == 0 /\ year %  
=> day <= 28)
```

# Invariant II

- ▶ All public methods of a class must preserve the class invariant.
- ▶ Class invariants can be incorporated into pre- and postconditions.

## Specification (of a method)

*Requires: Precondition and Invariant*

*Ensures: Postcondition and Invariant*

## Specification (of a constructor)

*Requires: Precondition*

*Ensures: Invariant*

## Further Elements of a Contract

Type signature (minimal contract)

Exceptions raised

Temporal properties

- ▶ the capacity of the table does not change over time
- ▶ a set that is only supposed to grow

# Testing vs. Verification

## TESTING

Goal: find evidence for **presence** of failures

Testing: execute a program with the intent of detecting failure

Testing cannot guarantee correctness, i.e., absence of failures

Related techniques: code reviews, program inspections

## VERIFICATION

Goal: find evidence for **absence** of failures

Verification guarantees correctness

Related techniques: code generation, program synthesis (from spec)

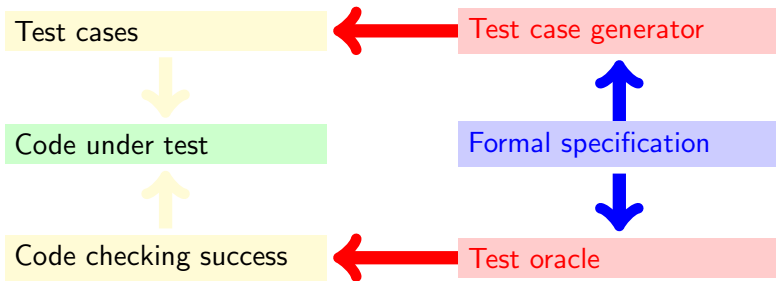
# Debugging: from Failures to Defects

- ▶ Both, testing and verification attempts exhibit **new** failures
- ▶ **Debugging** is a systematic process that finds and eliminates the defect that led to an observed failure
- ▶ Programs without **known** failures may still contain defects:
  - ▶ if they have not been verified
  - ▶ if they **have been** verified,  
but the failure is not covered by the specification

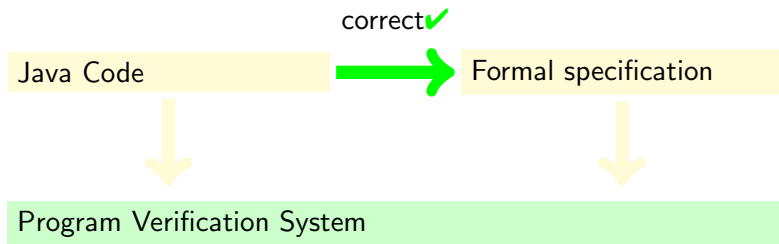
# Where Formalization Comes In

Testing is very expensive, even with tool support

30–80% of development time goes into testing



# Formal Verification of Program Correctness



Computer support essential for verification of real programs

```
synchronized java.lang.StringBuffer append(char c)
```

- ▶ ca. 15.000 proof steps
- ▶ ca. 200 case distinctions
- ▶ Two human interactions, ca. 1 minute computing time



# Tool Support is Essential

## Some Reasons for Using Tools

- ▶ Automate repetitive tasks
- ▶ Avoid typos, etc.
- ▶ Cope with large programs

## Tools Used

- ▶ Automated running of tests: `JUNIT`
- ▶ Debugging: `ECLIPSE` debugger