Software Engineering

Lecture 12: Testing and Debugging — Debugging

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Today's Topic

— Last Lecture —

- ✔ Bug tracking
- ✔ Program control — Design for Debugging
- ✔ Input simplification

Today's Topic

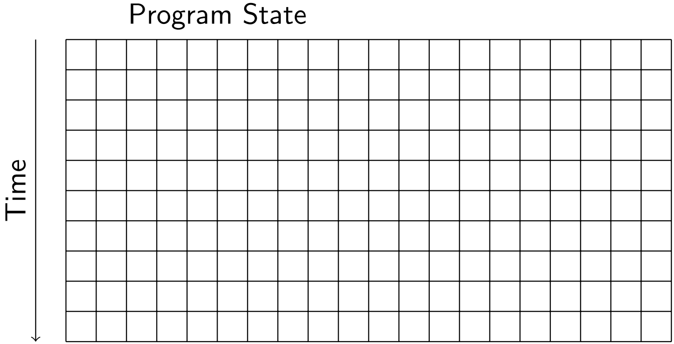
— Last Lecture —

- ✓ Bug tracking
- ✓ Program control — Design for Debugging
- ✓ Input simplification

— This Lecture —

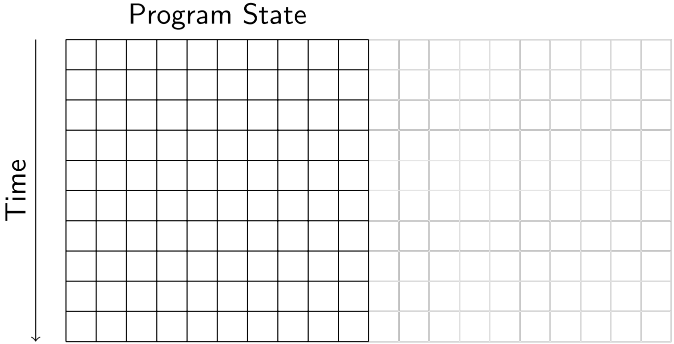
- ▶ Execution observation
  - ▶ With logging
  - ▶ Using debuggers
- ▶ Tracking causes and effects

The Main Steps in Systematic Debugging



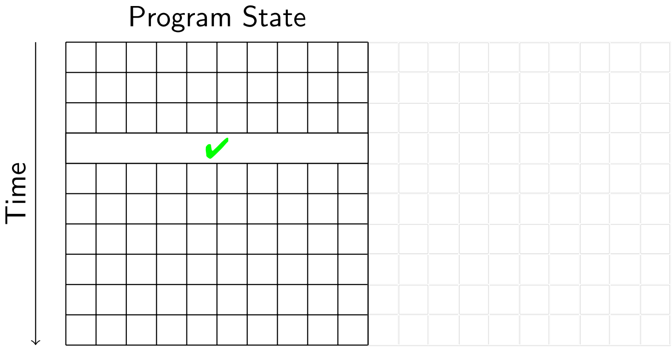
Reproduce failure with test input

The Main Steps in Systematic Debugging



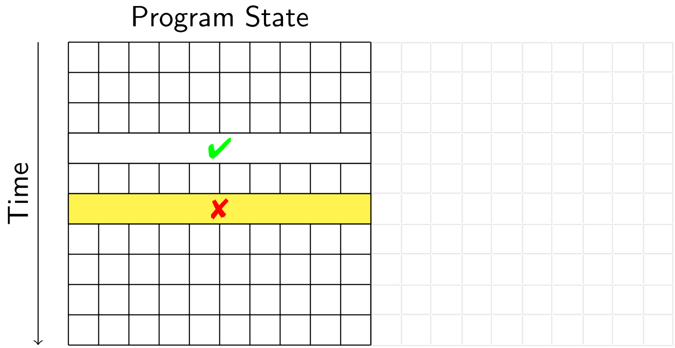
Reduction of failure-inducing problem

The Main Steps in Systematic Debugging



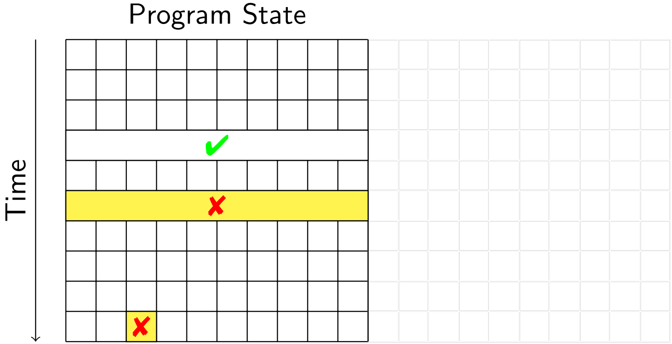
State known to be healthy

The Main Steps in Systematic Debugging



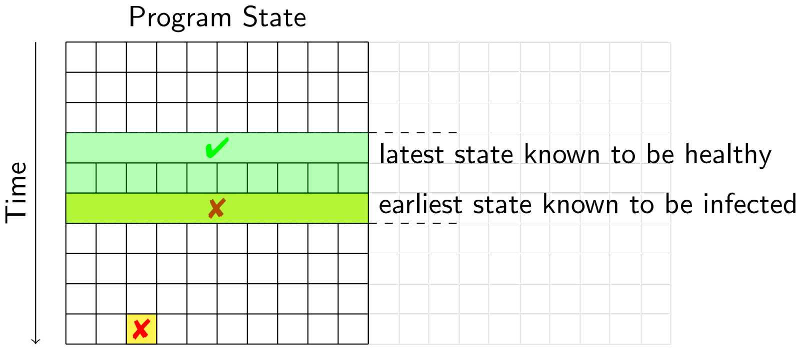
State known to be infected

The Main Steps in Systematic Debugging



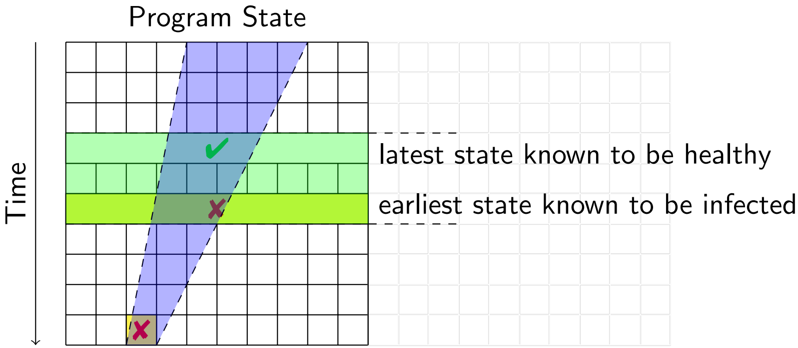
State where failure becomes observable

The Main Steps in Systematic Debugging



- Separate healthy from infected states

The Main Steps in Systematic Debugging



- Separate healthy from infected states
- Separate relevant from irrelevant states

Central Problem

How can we observe a program run?

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How can we observe a program run?

Challenges/Obstacles

- Observation of intermediate state not part of functionality
- Observation can change the behavior
- Narrowing down to relevant time/state sections

The Naive Approach: Print Logging

Println Debugging

Manually add print statements at code locations to be observed

```
System.out.println("size␣=␣"+ size);
```

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- ✓ Simple and easy
- ✓ Can use any output channel
- ✓ No tools or infrastructure needed, works on any platform

The Naive Approach: Print Logging

Println Debugging

Manually add print statements at code locations to be observed

```
System.out.println("size␣=␣"+ size);
```

- ✔ Simple and easy
- ✔ Can use any output channel
- ✔ No tools or infrastructure needed, works on any platform

- ✘ Code cluttering
- ✘ Output cluttering (at least need to use debug channel)
- ✘ Performance penalty, possibly changed behavior (timing, . . .)
- ✘ Buffered output lost on crash
- ✘ Source code required, recompilation necessary

Logging Frameworks

Example (Logging Framework for JAVA)

`java.util.logging`

Main principles of Java logging

- Each class can have its own `Logger` object
- Each logger is associated with a level and a handler
- Levels: `FINEST` < `FINER` < `FINE` < `CONFIG` < `INFO` < `WARNING` < `SEVERE`
- Handlers: `j.u.l.ConsoleHandler`, `j.u.l.FileHandler`
- Example: log `message` with `myLogger` and level `INFO`: `myLogger.info(Object message);`
- Logging can be controlled by program or properties file: which logger, level, filter, formatting, handler, etc.
- No recompilation necessary for reconfiguration

Evaluation of Logging Frameworks

- ✔ Output cluttering can be mastered
- ✔ Small performance overhead
- ✔ Exceptions are loggable
- ✔ Log complete up to crash
- ✔ Instrumented source code reconfigurable w/o recompilation

- ✘ Code cluttering — don't try to log everything!

Code cluttering avoidable with aspects, but also with Debuggers

What is a Debugger?

Basic Functionality of a Debugger

Execution Control	Stop execution on specified conditions: breakpoints
Interpretation	Step-wise execution of code
State Inspection	Observe value of variables and stack
State Change	Change state of stopped program

Historical term Debugger is misnomer as there are many debugging tools

What is a Debugger?

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Historical term Debugger is misnomer as there are many debugging tools

- Traditional debuggers (`gdb` for C) based on command line I/F
- We use the built-in GUI-based debugger of the ECLIPSE framework
  - Feel free to experiment with other debuggers!

Running Example

```
public static int search( int[] array,
                          int target ) {

  int low = 0;
  int high = array.length;
  int mid;
  while ( low <= high ) {
    mid = (low + high)/2;
    if ( target < array[ mid ] ) {
        high = mid - 1;
    } else if ( target > array[ mid ] ) {
        low = mid + 1;
    } else {
        return mid;
    }
  }
  return -1;
}
```

Eclipse Debugger

- Open directory `BinSearch`, create project `Search`
- Create/show run configuration `testBin1`
- Run `testBin1`
- Open Debugging view of project `Search`

Testing

Running a few test cases . . .

`search( {1,2}, 1 ) == 0` ✔

Testing

Running a few test cases . . .

search({1,2}, 1) == 0 ✓

search({1,2}, 2) == 1 ✓

Testing

Running a few test cases ...

`search( {1,2}, 1 ) == 0` ✔

`search( {1,2}, 2 ) == 1` ✔

`search( {1,2}, 4 )` throws

`ArrayIndexOutOfBoundsException:  3` ✘

Testing

Running a few test cases …

```
search( {1,2}, 1 ) == 0 ✔
search( {1,2}, 2 ) == 1 ✔
search( {1,2}, 4 ) throws
ArrayIndexOutOfBoundsException:  3 ✘
```

Example taken from a published JAVA text book :-(

Halting Program Execution

Breakpoint

A program location that, when it is reached, halts execution

Example (Setting Breakpoint)

In `search()` at loop, right-click, toggle breakpoint

Some remarks on breakpoints

- Set breakpoint at last statement where state is known to be healthy
- Formulate healthiness as an explicit hypothesis
- In Eclipse, not all lines can be breakpoints, because these are actually inserted into bytecode
- Remove breakpoints when no longer needed

Resuming Program Execution

Example (Execution Control Commands)

- Start debugging of run configuration `testBin1`
- Resume halts when breakpoint is reached in next loop execution
- Disable breakpoint for this session
- Resume executes now until end
- Remove from debug log (Remove All Terminated)
- Enable breakpoint again in Breakpoints window
- Close debugging perspective

Step-Wise Execution of Programs

Step-Wise Execution Commands

Step Into	Execute next statement, then halt
Step Over	Consider method call as one statement

Some remarks on step-wise execution

- Usually JAVA library methods stepped over
  - They should not contain defects
  - You probably don't have the source code
- To step over bigger chunks, change breakpoints, then resume

Inspecting the Program State

Inspection of state while program is halted

- Variables window
    - Unfold reference types
    - Pretty-printed in lower half of window
- Tooltips for variables in focus in editor window
- Recently changed variables are highlighted

Inspecting the Program State

Inspection of state while program is halted

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Example (Tracking `search()`)

- Start debugging at beginning of loop (`testBin2`)
- Step through one execution of loop body
- After first execution of loop body `low==high==2`
- Therefore, `mid==2`, but `array[2]` doesn't exist!
- If `target` is greater than all array elements, eventually `low==mid==array.length`

Changing the Program State

Hypothesis for Correct Value

Variable `high` should have value `array.length-1`

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Changing state while program is halted

- Right-click on identifier in Variables window, Change Value

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Changing state while program is halted

- Right-click on identifier in Variables window, Change Value

Example (Fixing the defect in the current run)

At start of second round of loop, set `high` to correct value `1`

Resuming execution now yields correct result

Watching States with Debuggers

Halting Execution upon Specific Conditions

Use Boolean Watch expression in conditional breakpoint

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Halting Execution upon Specific Conditions

Use Boolean Watch expression in conditional breakpoint

Example (Halting just before exception is thrown)

- From test run: argument `mid` of `array` is `2` at this point
- Create breakpoint at code position where evaluation takes place
- Add watch expression `mid==2` to breakpoint properties
- Disable breakpoint at start of loop
- Execution halts exactly when `mid==2` becomes `true`

Watching States with Debuggers

Halting Execution upon Specific Conditions

Use Boolean Watch expression in conditional breakpoint

Example (Halting just before exception is thrown)

- From test run: argument `mid` of `array` is 2 at this point
- Create breakpoint at code position where evaluation takes place
- Add watch expression `mid==2` to breakpoint properties
- Disable breakpoint at start of loop
- Execution halts exactly when `mid==2` becomes `true`

Hints on watch expressions

- Make sure scope of variables in watch expressions is big enough

Evaluation of Debuggers

- ✔ Code cluttering completely avoided
- ✔ Prudent usage of breakpoints/watches reduces states to be inspected
- ✔ Full control over all execution aspects

- ✘ Debuggers are interactive tools, re-use difficult
- ✘ Performance can degrade, disable unused watches
- ✘ Inspection of reference types (lists, etc.) is tedious

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Important Lessons

- ▶ Both, logging and debuggers are necessary and complementary
- ▶ Need visualization tools to render complex data structures
- ▶ Minimal/small input, localisation of unit is important

Tracking Causes and Effects

Determine defect that is origin of failure

Fundamental problem

Program executes forward, but need to reason backwards from failure

Example

In `search()` the failure was caused by wrong value `mid`,
but the real culprit was `high`

Effects of Statements

Fundamental ways how statements may affect each other

Write	Change the program state Assign a new value to a variable read by another statement
Control	Change the program counter Determine which statement is executed next

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Statements with Write Effect (in Java)

- Assignments
- I/O, because it affects buffer content
- new(), because object initialisation writes to fields

Effects of Statements

Fundamental ways how statements may affect each other

Write	Change the program state Assign a new value to a variable read by another statement
Control	Change the program counter Determine which statement is executed next

Statements with Control Effect (in JAVA)

- Conditionals, switches
- Loops: determine whether their body is executed
- Dynamic method calls: implicit case distinction on implementations
- Abrupt termination statements: `break`, `return`
- Exceptions: potentially at each object or array access!

Statement Dependencies

Definition (Data Dependency)

Statement `B` is data dependent on statement `A` iff

1. `A` writes to a variable `v` that is read by `B` and
2. There is at least one execution path between `A` and `B` in which `v` is not written to

"The outcome of `A` can directly influence a variable read in `B`"

Statement Dependencies

Definition (Control Dependency)

Statement `B` is control dependent on statement `A` iff

- There is an execution path from `A` to `B` such that:
  For all statements `S` $\neq$`A` on the path, all execution paths from `S` to the method exit pass through `B`
  and
- There is an execution path from `A` to the method exit that does **not** pass through `B`

“The outcome of `A` can influence whether `B` is executed”

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

Example

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int low = 0;
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  }
}
return -1;
```

`mid` is data-dependent on this statement

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

`mid` is control-dependent on the while statement

Computing Backward Dependencies

Definition (Backward Dependency)

Statement `B` is backward dependent on statement `A` iff

There is a sequence of statements $\mathtt{A} = \mathtt{A}_1, \mathtt{A}_2, \ldots, \mathtt{A}_\mathtt{n} = \mathtt{B}$ such that:

1. for all *i*, $\mathtt{A}_{i+1}$ is either control dependent or data dependent on $\mathtt{A}_i$
2. there is at least one *i* with $\mathtt{A}_{i+1}$ being data dependent on $\mathtt{A}_i$

"The outcome of `A` can influence the program state in `B`"

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
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`mid` is backward-dependent on data- and control- dependent statement

Example

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int low = 0;
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while ( low <= high ) {
  mid = (low + high)/2;
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`mid` is backward-dependent on data- and control- dependent statement

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  } else {
      return mid;
  }
}
return -1;
```

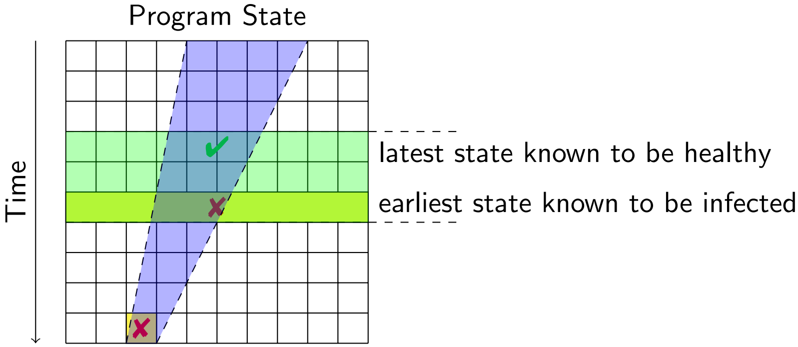
Backward-dependent statements for first execution of loop body

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

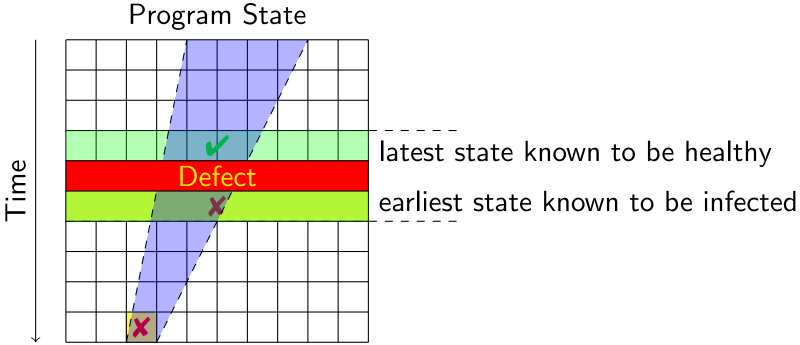
Backward-dependent statements for repeated execution of loop body

Systematic Discovery of Defects



- Separate healthy from infected states
- Separate relevant from irrelevant states

Systematic Discovery of Defects



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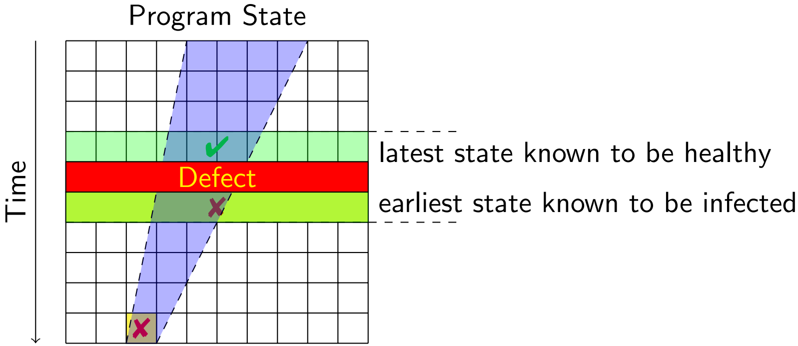
Systematic Discovery of Defects



- Separate healthy from infected states
- Separate relevant from irrelevant states
- Compute backward-dependent statements from infected locations

Algorithm: Systematic Discovery of Defects

Invariant: $\mathcal{I}$ is a set of locations (variable set V and statement S) such that each $v \in V$ is infected after executing S.

```
1. Initialize $\mathcal{I} := \{\text{infected location reported by failure}\}$
2. Choose a current, infected location $L = (V, S) \in \mathcal{I}$
3. Let $\mathcal{I} := \mathcal{I} \setminus \{L\}$
4. Let $\mathcal{C} := \emptyset$ accumulate a set of candidates
5. For each statement $S'$ that may contain origin of defect:
   $S$ backwards depends on $S'$ in one step in execution path
   5.1 Let $\mathcal{M}$ be the set of variables that is written in $S'$ and infected
   5.2 If $\mathcal{M} \neq \emptyset$ let $\mathcal{C} := \mathcal{C} \cup \{(\mathcal{M}, S')\}$
6. If $\mathcal{C} \neq \emptyset$ (there are infected predecessors):
   6.1 Let $\mathcal{I} := \mathcal{I} \cup \mathcal{C}$
   6.2 Goto 2.
7. $L$ depends only on healthy locations, it must be the infection site!
```

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

`mid` is infected, `mid==low==high==2`

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

Look for origins of `low` and `high`

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

`low` was changed in previous loop execution, value `low==1` seems healthy

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

`high`==2 set at start (if-branch not taken when target not found), infecte

Example

```
int low = 0;
int high = array.length;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

`high` does not depend on any other location—found infection site!

Example

```
int low = 0;
int high = array.length - 1;
int mid;
while ( low <= high ) {
  mid = (low + high)/2;
  if ( target < array[ mid ] ) {
      high = mid - 1;
  } else if ( target > array[ mid ] ) {
      low = mid + 1;
  } else {
      return mid;
  }
}
return -1;
```

Fixed defect

After Fixing the Defect

- Failures that exhibited a defect become new test cases after the fix
  - used for regression testing
- Use existing unit test cases to
  - test a suspected method in isolation
  - make sure that your bug fix did not introduce new bugs
  - exclude wrong hypotheses about the defect

Open Questions

1. How is evaluation of test runs related to specification?
   So far: wrote oracle program or evaluated interactively
   How to check automatically whether test outcome conforms to spec?
2. It is tedious to write test cases by hand
   Easy to forget cases
   JAVA: aliasing, run-time exceptions
3. When does a program have no more bugs?
   How to prove correctness without executing ∞ many paths?

Literature for this Lecture

Essential

Zeller	Why Programs Fail: A Guide to Systematic Debugging, Morgan Kaufmann, 2005 Chapters 7, 8, 9